## 1. Introduction

- course introduction
- optimization examples
- vectors and matrices


## Mathematical optimization

## (mathematical) optimization problem

$$
\begin{align*}
\operatorname{minimize} \text { (or maximize) } & f(\boldsymbol{x}) \\
\text { subject to } & g_{i}(\boldsymbol{x}) \leq 0, i=1, \ldots, m  \tag{1.1}\\
& h_{j}(\boldsymbol{x})=0, j=1, \ldots, p
\end{align*}
$$

- optimization variable: $\boldsymbol{x}=\left(x_{1}, \ldots, x_{n}\right) \in \mathbb{R}^{n}$
- objective function: $f: \mathbb{R}^{n} \rightarrow \mathbb{R}$
- inequality constraints functions: $g_{i}: \mathbb{R}^{n} \rightarrow \mathbb{R}$
- equality constraints functions: $h_{j}: \mathbb{R}^{n} \rightarrow \mathbb{R}$

Optimal point or solution: a point $\boldsymbol{x}^{\star}$ is an optimal point or solution to problem (1.1) if it attains the smallest (largest) objective value among all points that satisfy the constraints

## Applications

## Applications

- electrical network design
- min-weight aircraft and aerospace design
- optimal space vehicle trajectories
- cost-efficient civil structure design (e.g., bridges, dams)
- min-cost material handling equipment design (e.g., trucks, cranes)
- maximizing profit investment strategies
- optimal data model fitting

Modeling: the process of identifying the objective, constraints, and variables of a given problem is called modeling

## Finding best decision

- the variable, $\boldsymbol{x}$, symbolizes specific actions such as:
- trades in a portfolio
- adjustments to airplane control surfaces
- task scheduling or assignment
- resource allocation decisions
- transmitted signal...
- constraint functions limit the action or set conditions on outcome:
- physical or technical limits
- resource budgets
- design requirements that need be satisfied...
- objective represents some criteria, we want to minimize:
- total cost
- deviation from desired outcome (error)
- consumption of fuel
- risk...


## Linear and nonlinear optimization

an optimization problem is called linear program if it has the form

$$
\begin{aligned}
\text { minimize (or maximize) } & \sum_{i=1}^{n} c_{i} x_{i} \\
\text { subject to } & \sum_{j=1}^{n} a_{i j} x_{j} \leq b_{i}, \quad i=1, \ldots, m \\
& \sum_{j=1}^{n} g_{i j} x_{j}=h_{i}, \quad i=1, \ldots, p
\end{aligned}
$$

- $\left\{c_{i}, a_{i j}, g_{i j}, h_{i}, b_{i}\right\}$ are given coefficients
- the objective and constraint functions are "linear"

Nonlinear program: an optimization problem is called nonlinear program if it is not a linear program

## Other optimization classes:

- Unconstrained optimization: no constraints, i.e., $h_{j}(\boldsymbol{x})=g_{i}(\boldsymbol{x})=0$
- Discrete optimization: variables take only discrete or integer values
- Integer linear program: a discrete optimization with linear objective and constraints
- Mixed integer optimization: variables can be both integer and continuous
(Note: this course focuses solely on optimization with continuous variables)


## Solving optimization problems

- various methods exist to solve optimization problems
- the chosen method will typically depend on several factors (e.g., problem class and structure)
- solutions guide decision-makers, who oversee, validate, and adjust the approach or problem as required.


## Course topics

## General course topics

- unconstrained and constrained optimization: optimality conditions
- convex optimization and duality
- solution methods: unconstrained and constrained
- modeling and applications in optimization


## Prerequisites

- solid foundation in linear algebra and calculus
- programming in MATLAB: prior experience not mandatory, but self-study is expected in the course's early stages


## Course objectives

- learn the basic mathematical theory of nonlinear optimization (and convex optimization) and their applications in practice
- learn and implement basic (some advanced) optimization methods
- develop the ability to identify optimization problem types and select appropriate solution methods
- equip yourself with optimization knowledge needed for research and real-world applications


## Course information

Course materials: lecture slides and other course material will be posted on Moodle

## Grading

- homework (20\%)
- two midterm exams (50\%)
- project (30\%)
(these weights are approximate; we reserve the right to change them later)
refer to the syllabus on the Moodle course website for more information, such as course references, office hours, class policy, exam dates, etc.


## Al tools policy

- unauthorized use of Al tools, like ChatGPT, is treated as plagiarism
- Al is an aid, not a substitute for genuine understanding; reliance solely on Al without understanding can result in penalties
- suspected misuse of AI may lead to oral exams or alternative assessments


## Outline

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## Maximum power transfer



- voltage source: $V$ (in volts)
- line resistor: $R$ (given value)
- objective: Determine $R_{L}$ to maximize power to it
using circuit analysis, the power delivered to $R_{L}$ is $p\left(R_{L}\right)=i^{2} R_{L}$ and $i=V /\left(R+R_{L}\right)$; hence, we can formulate the problem as

$$
\text { maximize } \frac{V^{2} x}{(R+x)^{2}}
$$

with variable $x=R_{L}$; this is an unconstrained nonlinear program

## Battery charging

an electric circuit is designed to use a 30 V source to charge $10 \mathrm{~V}, 6 \mathrm{~V}$, and 20 V batteries


- physical constraints limit the currents $i_{1}, i_{2}, i_{3}, i_{4}$, and $i_{5}$ to a maximum of $4 \mathrm{~A}, 3 \mathrm{~A}, 3 \mathrm{~A}, 2 \mathrm{~A}$, and 2 A
- the batteries must not be discharged; that is, the currents $i_{1}, i_{2}, i_{3}, i_{4}$, and $i_{5}$ must not be negative
- we wish to find the values of the currents $i_{1}, i_{2}, \ldots, i_{5}$ such that the total power transferred to the batteries is maximized
using circuit analysis, the problem can be modeled as the linear program:

$$
\begin{array}{cl}
\operatorname{maximize} & 10 i_{2}+6 i_{4}+20 i_{5} \\
\text { subject to } & i_{1}=i_{2}+i_{3} \\
& i_{3}=i_{4}+i_{5} \\
& i_{1} \leq 4 \\
& i_{2} \leq 3 \\
& i_{3} \leq 3 \\
& i_{4} \leq 2 \\
& i_{5} \leq 2 \\
& i_{1}, i_{2}, i_{3}, i_{4}, i_{5} \geq 0
\end{array}
$$

once the currents are found, we can find the resistors $R_{1}, \ldots, R_{5}$ that draw such currents using Ohm's and Kirchhoff's laws

## Concrete mixture

Concrete type 1:

- Cost: \$5/lb
- Cement: 30\%
- Gravel: $40 \%$
- Sand: 30\%

Concrete type 2:

- Cost: \$1/lb
- Cement: 10\%
- Gravel: 20\%
- Sand: 70\%

Formulate a mixture with at least: 5 kg of Cement, 3 kg of Gravel, 4 kg of Sand, while minimizing cost

## Problem formulation:

$$
\begin{array}{ll}
\operatorname{minimize} & 5 x_{1}+x_{2} \\
\text { subject to } & 0.3 x_{1}+0.1 x_{2} \geq 5 \\
& 0.4 x_{1}+0.2 x_{2} \geq 3 \\
& 0.3 x_{1}+0.7 x_{2} \geq 4 \\
& x_{1} \geq 0, x_{2} \geq 0
\end{array}
$$

variables $x_{i}$ represent the weight of concrete $i$ that we want to buy

## Knapsack problem

## Description

- given $n$ items, each with a weight and value
- determine the number of each item to include, maximizing total value while ensuring the total weight doesn't exceed a set limit

Investment example:

- Aim: invest among $n$ opportunities
- Budget: at most $d$ dollars
- $i$ th investment:
- cost: $c_{i}$ dollars
- expected profit: $p_{i}$
- available units: $b_{i}$
how many items of each type should be bought to maximize the expected profit?
problem can be formulated as

$$
\begin{aligned}
\text { maximize } & \sum_{i=1}^{n} p_{i} x_{i} \\
\text { subject to } & \sum_{i=1}^{n} c_{i} x_{i} \leq d, \text { (total cost } \leq \text { available amount), } \\
& x_{i} \in\left\{0,1,2, \ldots, b_{i}\right\}, \quad i=1, \ldots, n
\end{aligned}
$$

- the problem is an integer optimization problem since the variables are restricted to be integers
- specifically, it is an integer linear program since the objective and constraints are "linear" and the variables are integer


## Facility placement: Fermat-Weber problem

consider locations of some facilities represented by points:
$\left(a_{1}, b_{1}\right), \ldots,\left(a_{m}, b_{m}\right)$ in 2D space

- Goal: Determine the optimal location, $\boldsymbol{x}=\left(x_{1}, x_{2}\right)$, of a distribution center to minimize the total daily distance.
- distance between the center, $\boldsymbol{x}$, and a facility, $\left(a_{i}, b_{i}\right)$, is:

$$
d_{i}=\sqrt{\left(x_{1}-a_{i}\right)^{2}+\left(x_{2}-b_{i}\right)^{2}}
$$

the problem can be formulated as

$$
\operatorname{minimize} \quad \sum_{i=1}^{m} w_{i} \sqrt{\left(x_{1}-a_{i}\right)^{2}+\left(x_{2}-b_{i}\right)^{2}}
$$

- $w_{i}$ is the weight associated with distance $d_{i}$ (e.g., higher values of $w_{i}$ for areas with more traffic)
- this problem is known as the Fermat-Weber problem


## Electrical wires connections

four buildings are to be connected by electrical wires


- central joining point: $\left(x_{0}, y_{0}\right)$
- each building $i$ connects at $\left(x_{i}, y_{i}\right)$ with wire length $d_{i}$
- Objective: find the positions $\left(x_{i}, y_{i}\right)$ that minimize the total length of wires used
- building 1 (circular): center $(1,4)$, radius 2
- building 2 (circular): center $(9,5)$, radius 1
- building 3 (square): center $(3,-2)$, side length 2
- building 4 (rectangle): center ( 7,0 ), height 4 , width 2


## Problem Formulation:

$$
\begin{array}{ll}
\text { minimize } & \sum_{i=1}^{4} \sqrt{\left(x_{i}-x_{0}\right)^{2}+\left(y_{i}-y_{0}\right)^{2}} \\
\text { subject to } & \left(x_{1}-1\right)^{2}+\left(y_{1}-4\right)^{2} \leq 4 \\
& \left(x_{2}-9\right)^{2}+\left(y_{2}-5\right)^{2} \leq 1 \\
& 2 \leq x_{3} \leq 4 \\
& -3 \leq y_{3} \leq-1 \\
& 6 \leq x_{4} \leq 8 \\
& -2 \leq y_{4} \leq 2
\end{array}
$$

with variables $\left(x_{i}, y_{i}\right)(i=0,1, \ldots, 4)$

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## Vectors

column vector of size $n$ (or $n$-vector):

$$
\boldsymbol{a}=\left[\begin{array}{c}
a_{1} \\
a_{2} \\
\vdots \\
a_{n}
\end{array}\right]
$$

- $a_{i}: i$ th entry (also termed as element, coefficient, or component)
- alternate notation: $\boldsymbol{a}=\left(a_{1}, \ldots, a_{n}\right)$
- set of real $n$-vectors: $\mathbb{R}^{n}$
- transpose results in a row vector: $\boldsymbol{a}^{T}=\left[\begin{array}{lll}a_{1} & \cdots & a_{n}\end{array}\right]$

Special vectors (dimension determined from context)

- zero vector: $\mathbf{0}=(0,0, \ldots, 0)$
- one vector: $\mathbf{1}=(1,1, \ldots, 1)$
- unit vectors: $\boldsymbol{a}=\boldsymbol{e}_{i}$ has entry $a_{i}=1$; all other entries are equal to zero


## Vectors operations and properties

- scalar vector multiplication: $\boldsymbol{x}=\alpha \boldsymbol{a}$, where $x_{i}=\alpha a_{i}$
- vector addition: $\boldsymbol{x}=\boldsymbol{a}+\boldsymbol{b}$, where $x_{i}=a_{i}+b_{i}$


## Vector addition properties

- commutative: $\boldsymbol{a}+\boldsymbol{b}=\boldsymbol{b}+\boldsymbol{a}$
- associative: $(\boldsymbol{a}+\boldsymbol{b})+\boldsymbol{c}=\boldsymbol{a}+(\boldsymbol{b}+\boldsymbol{c})$

Scalar-vector multiplication properties

- distributive: for any real scalars $\alpha$ and $\beta$

$$
\begin{aligned}
& \alpha(\boldsymbol{a}+\boldsymbol{b})=\alpha \boldsymbol{a}+\alpha \boldsymbol{b} \\
& (\alpha+\beta) \boldsymbol{a}=\alpha \boldsymbol{a}+\beta \boldsymbol{a}
\end{aligned}
$$

- associative: $\alpha(\beta \boldsymbol{a})=(\alpha \beta) \boldsymbol{a}$, which we write as $\alpha \beta \boldsymbol{a}$


## Geometric interpretation: displacement



- Left: the displacement $\boldsymbol{a}+\boldsymbol{b}$
- Right: the displacement $\boldsymbol{b}+\boldsymbol{a}$


## Geometric interpretation: position



- the point $\boldsymbol{b}+\boldsymbol{a}$ is the position of the point represented by $\boldsymbol{b}$ displaced by the displacement represented by $a$
- the vector $\boldsymbol{p}-\boldsymbol{q}$ represents the displacement from the point represented by $\boldsymbol{q}$ to the point represented by $\boldsymbol{p}$


## Line segment

any point on the line passing through distinct $\boldsymbol{a}$ and $\boldsymbol{b}$ can be written as

$$
\boldsymbol{c}=\theta \boldsymbol{a}+(1-\theta) \boldsymbol{b}
$$

- $\theta$ is a scalar
- for $0 \leq \theta \leq 1$, point $\boldsymbol{c}$ lie on the segment between $\boldsymbol{a}$ and $\boldsymbol{b}$



## Euclidean inner product and norm

(euclidean) inner product of two $n$-vectors:

$$
\boldsymbol{a}^{T} \boldsymbol{b}=\sum_{i=1}^{n} a_{i} b_{i}=a_{1} b_{1}+\cdots+a_{n} b_{n}
$$

$\boldsymbol{a}$ and $\boldsymbol{b}$ are said to be orthogonal if $\boldsymbol{a}^{T} \boldsymbol{b}=0$

## Properties

- nonnegativity: $\boldsymbol{a}^{T} \boldsymbol{a} \geq 0$ and $\boldsymbol{a}^{T} \boldsymbol{a}=0$ if and only if $\boldsymbol{a}=\mathbf{0}$
- symmetry: $\boldsymbol{a}^{T} \boldsymbol{b}=\boldsymbol{b}^{T} \boldsymbol{a}$
- additivity: $(\boldsymbol{a}+\boldsymbol{b})^{T} \boldsymbol{c}=\boldsymbol{a}^{T} \boldsymbol{c}+\boldsymbol{b}^{T} \boldsymbol{c}$
- homogeneity: $(\alpha \boldsymbol{a})^{T} \boldsymbol{b}=\alpha\left(\boldsymbol{a}^{T} \boldsymbol{b}\right)$


## Euclidean norm

$$
\|\boldsymbol{a}\|=\sqrt{a_{1}^{2}+a_{2}^{2}+\cdots+a_{n}^{2}}=\sqrt{\boldsymbol{a}^{T} \boldsymbol{a}}
$$

## Cauchy-Schwarz inequality

$$
\left|\boldsymbol{a}^{T} \boldsymbol{b}\right| \leq\|\boldsymbol{a}\|\|\boldsymbol{b}\|
$$

## Inner product examples

- unit vector: inner product $\boldsymbol{e}_{i}^{T} \boldsymbol{a}=a_{i}$ picks $i$ th element of $\boldsymbol{a}$
- sum: $\mathbf{1}^{T} \boldsymbol{a}=a_{1}+a_{2}+\cdots+a_{n}$
- average: $(\mathbf{1} / n)^{T} \boldsymbol{a}=\frac{a_{1}+a_{2}+\cdots+a_{n}}{n}$
- sum of squares: $\boldsymbol{a}^{T} \boldsymbol{a}=a_{1}^{2}+a_{2}^{2}+\cdots+a_{n}^{2}=\|\boldsymbol{a}\|^{2}$


## Norm of block vectors

let $\boldsymbol{a}, \boldsymbol{b}$ and $\boldsymbol{c}$, be any real $n$-vectors and define the block vector:

$$
d=(a, b, c)=\left[\begin{array}{l}
a \\
b \\
c
\end{array}\right]
$$

- the euclidean norm squared of block vector $\boldsymbol{d}$ is

$$
\|\boldsymbol{d}\|^{2}=\|\boldsymbol{a}\|^{2}+\|\boldsymbol{b}\|^{2}+\|\boldsymbol{c}\|^{2}
$$

- the norm of $\boldsymbol{d}=(\boldsymbol{a}, \boldsymbol{b}, \boldsymbol{c})$ can be expressed as:

$$
\|\boldsymbol{d}\|=\sqrt{\|\boldsymbol{a}\|^{2}+\|\boldsymbol{b}\|^{2}+\|\boldsymbol{c}\|^{2}}=\|(\|\boldsymbol{a}\|,\|\boldsymbol{b}\|,\|\boldsymbol{c}\|)\|
$$

## Matrices

matrix of size (dimension) $m \times n$ :

$$
A=\left[\begin{array}{cccc}
a_{11} & a_{12} & \cdots & a_{1 n} \\
a_{21} & a_{22} & \cdots & a_{2 n} \\
\vdots & \vdots & & \vdots \\
a_{m 1} & a_{m 2} & \cdots & a_{m n}
\end{array}\right]
$$

- $a_{i j}$ is the $i, j$ element (or entry or coefficient)
- set of real $m \times n$ matrices is denoted by $\mathbb{R}^{m \times n}$
- transpose of $A$ : $A^{T}$ is an $n \times m$ matrix with $i, j$ entry equal to $a_{j i}$
- a matrix is square if $m=n$


## Special matrices:

- zero matrix: $A=0$ is a matrix with zeros entries $a_{i j}=0$
- diagonal matrix: square matrix $A=\operatorname{diag}\left(a_{1}, \ldots, a_{n}\right)$ with diagonal entries $a_{i i}=a_{i}$ and off-diagonal $a_{i j}=0$
- identity matrix: diagonal matrix $A=I, m=n$ with diagonal entries $a_{i i}=1$ and $a_{i j}=0$
- symmetric matrix: square matrix $A$ with $A=A^{T}$


## Columns and rows of a matrix

an $m \times n$ matrix

$$
A=\left[\begin{array}{cccc}
a_{11} & a_{12} & \cdots & a_{1 n} \\
a_{21} & a_{22} & \cdots & a_{2 n} \\
\vdots & \vdots & & \vdots \\
a_{m 1} & a_{m 2} & \cdots & a_{m n}
\end{array}\right]
$$

- $\boldsymbol{a}_{j}=\left(a_{1 j}, \ldots, a_{m j}\right)$ is the $j$ th column of $A$
- $\hat{\boldsymbol{a}}_{i}^{T}=\left[\begin{array}{lll}a_{i 1} & \cdots & a_{i n}\end{array}\right]$ is the $i$ th row of $A$
- the matrix $A$ can represented in terms of its columns or row as

$$
A=\left[\begin{array}{llll}
\boldsymbol{a}_{1} & \boldsymbol{a}_{2} & \cdots & \boldsymbol{a}_{n}
\end{array}\right] \quad \text { or } \quad A=\left[\begin{array}{c}
\hat{\boldsymbol{a}}_{1}^{T} \\
\hat{\boldsymbol{a}}_{2}^{T} \\
\vdots \\
\hat{\boldsymbol{a}}_{m}^{T}
\end{array}\right]
$$

## Block matrices

matrices can be partitioned into submatrices; for example,

$$
A=\left[\begin{array}{ll}
B & C \\
D & E
\end{array}\right]
$$

denote a matrix where $B, C, D$, and $E$ are matrices themselves

- $B, C, D$, and $E$ are termed as blocks or submatrices
- blocks are identified by their row and column indices, for example, $C$ is the $(1,2)$ block of $A$
- block matrices must have compatible dimensions


## Example

$$
B=\left[\begin{array}{lll}
1 & 1 & 1
\end{array}\right], \quad C=[-1], \quad D=\left[\begin{array}{lll}
2 & 2 & 2 \\
3 & 3 & 3
\end{array}\right], \quad E=\left[\begin{array}{l}
4 \\
4
\end{array}\right]
$$

we have

$$
A=\left[\begin{array}{ll}
B & C \\
D & E
\end{array}\right]=\left[\begin{array}{rrrr}
1 & 1 & 1 & -1 \\
2 & 2 & 2 & 4 \\
3 & 3 & 3 & 4
\end{array}\right]
$$

## Matrix operations

- scalar matrix multiplication: $X=\alpha A, x_{i j}=\alpha a_{i j}$
- matrix addition: $X=A+B, x_{i j}=a_{i j}+b_{i j}$

Matrix-vector product: multiplication of a matrix $A \in \mathbb{R}^{m \times n}$ with a vector of compatible size $\boldsymbol{x} \in \mathbb{R}^{n}$ :

$$
\boldsymbol{y}=A \boldsymbol{x}, \quad y_{i}=\sum_{j=1}^{n} a_{i j} x_{j}, \quad i=1, \ldots, m
$$

Matrix-matrix product: product of a matrix with a matrix $A \in \mathbb{R}^{m \times p}$ of compatible size $B \in \mathbb{R}^{p \times n}, C=A B$ :

$$
c_{i j}=\sum_{\ell=1}^{p} a_{i \ell} b_{\ell j}=a_{i 1} b_{1 j}+\cdots+a_{i p} b_{p j}
$$

$i=1, \ldots, m, j=1, \ldots, n$

## Matrix operations properties

## Matrix addition properties

- commutativity: $A+B=B+A$
- associativity: $(A+B)+C=A+(B+C)$
- addition with zero matrix: $A+0=0+A=A$
- transpose of sum: $(A+B)^{T}=A^{T}+B^{T}$


## Matrix-vector multiplication properties

- distributive: $A(\boldsymbol{u}+\boldsymbol{v})=A \boldsymbol{u}+A \boldsymbol{v}$ and $(A+B) \boldsymbol{u}=A \boldsymbol{u}+B \boldsymbol{u}$ where $\boldsymbol{u}, \boldsymbol{v}$ are vectors and $A, B$ are matrices
- homogeneity: $(\alpha A) \boldsymbol{u}=\alpha(A \boldsymbol{u})=A(\alpha \boldsymbol{u})$, which we write as $\alpha A \boldsymbol{u}$


## Matrix multiplication properties

- associativity: $(A B) C=A(B C)$, which we write it as $A B C$
- distributivity with addition

$$
A(B+C)=A B+A C, \quad(A+B) C=A C+B C
$$

- transpose of product: $(A B)^{T}=B^{T} A^{T}$
- for scalars $\alpha$ and $\beta$, we have

$$
\begin{aligned}
\alpha(A B) & =(\alpha A) B=A(\alpha B) \\
(\alpha A)^{T} & =\alpha A^{T} \\
(\alpha \beta) A & =\alpha(\beta A) \\
(\alpha+\beta) A & =\alpha A+\beta A
\end{aligned}
$$

## Matrix trace

the trace of an $n \times n$ matrix $A$ is the sum of its diagonal elements:

$$
\operatorname{Tr}(A)=\sum_{i=1}^{n} a_{i i}
$$

some properties of the trace are:

- $\operatorname{Tr}(A)=\operatorname{Tr}\left(A^{T}\right)$
- if $A$ is an $m \times n$ matrix and $B$ is an $n \times m$ matrix, then

$$
\operatorname{Tr}(A B)=\operatorname{Tr}(B A)
$$

## Functions

- $f: \mathcal{X} \rightarrow \mathcal{Y}$ to denote a function $f$ that maps an element from the set $\mathcal{X}$ into the set $\mathcal{Y}$
- $f: \mathbb{R}^{n} \rightarrow \mathbb{R}^{m}$ means that $f$ maps a real $n$-vector to a real $m$-vector:

$$
f(\boldsymbol{x})=\left[\begin{array}{c}
f_{1}(\boldsymbol{x}) \\
\vdots \\
f_{m}(\boldsymbol{x})
\end{array}\right]
$$

where the entry $f_{i}: \mathbb{R}^{n} \rightarrow \mathbb{R}$ is itself a scalar-valued function of $\boldsymbol{x}$

Function domain: the domain of function $f$, symbolized by $\operatorname{dom} f \subseteq \mathcal{X}$, represents the set of points where $f$ is defined and finite; for example, the functions

$$
f_{1}(x)=\left\{\begin{array}{ll}
1 / x & \text { if } x \neq 0 \\
\infty & \text { otherwise }
\end{array}, \quad f_{2}(x)= \begin{cases}1 / x & \text { if } x>0 \\
\infty & \text { otherwise }\end{cases}\right.
$$

are different since they have different domains

## Examples

defined everywhere (dom $f=\mathbb{R}^{n}$ )

- $f: \mathbb{R} \rightarrow \mathbb{R}: f(x)=x^{2}+x+1$ maps a scalar $x$ to a scalar $f(x)$
- $f: \mathbb{R}^{3} \rightarrow \mathbb{R}: f\left(x_{1}, x_{2}, x_{3}\right)=x_{1}+x_{2}+x_{3}^{2}$
- $f: \mathbb{R}^{n} \rightarrow \mathbb{R}^{m}: f(\boldsymbol{x})=A \boldsymbol{x}$ where $\boldsymbol{x} \in \mathbb{R}^{n}$ and $A$ is an $m \times n$ matrix
- $f: \mathbb{R}^{2} \rightarrow \mathbb{R}^{3}: f\left(x_{1}, x_{2}\right)=\left(x_{1}, x_{2}, x_{1}+x_{2}^{2}\right)$


## undefined everywhere

- $f(x)=\log x(f: \mathbb{R} \rightarrow \mathbb{R})$ is a function that takes a real number and outputs a real number and it is valid only for $x>0$, hence $\operatorname{dom} f=\{x \mid x>0\}$
- $f\left(x_{1}, x_{2}\right)=x_{1} / x_{1}+x_{2}\left(f: \mathbb{R}^{2} \rightarrow \mathbb{R}\right)$ with $\operatorname{dom} f=\left\{\boldsymbol{x} \mid x_{1}+x_{2} \neq 0\right\}$ where $\boldsymbol{x}=\left(x_{1}, x_{2}\right)$


## Linear functions

a function $f: \mathbb{R}^{n} \rightarrow \mathbb{R}^{m}$ is linear if it satisfies the superposition property:

$$
f(\alpha \boldsymbol{x}+\beta \boldsymbol{y})=\alpha f(\boldsymbol{x})+\beta f(\boldsymbol{y})
$$

for any an $n$-vectors $\boldsymbol{x}, \boldsymbol{y}$ and any scalars $\alpha, \beta$

- a linear function $f: \mathbb{R}^{n} \rightarrow \mathbb{R}$ can always be expressed as: $f(\boldsymbol{x})=\boldsymbol{a}^{T} \boldsymbol{x}$ for some $n$-vector $\boldsymbol{a}$.
- similarly, a linear function $f: \mathbb{R}^{n} \rightarrow \mathbb{R}^{m}$ can be represented as: $f(\boldsymbol{x})=A \boldsymbol{x}$ for some $m \times n$ matrix $A$
- to see this, using the linear property of $f$, we have:

$$
f(\boldsymbol{x})=f\left(x_{1} \boldsymbol{e}_{1}+\cdots+x_{n} \boldsymbol{e}_{n}\right)=x_{1} f\left(\boldsymbol{e}_{1}\right)+\cdots+x_{n} f\left(\boldsymbol{e}_{n}\right)=A \boldsymbol{x}
$$

where the matrix $A$ has columns:

$$
A=\left[f\left(\boldsymbol{e}_{1}\right) \cdots f\left(\boldsymbol{e}_{n}\right)\right]
$$

## Affine functions

a function $f: \mathbb{R}^{n} \rightarrow \mathbb{R}^{m}$ is affine if it can be expressed as

$$
f(\boldsymbol{x})=A \boldsymbol{x}+\boldsymbol{b}
$$

for some $A \in \mathbb{R}^{m \times n}$ and $\boldsymbol{b} \in \mathbb{R}^{m}$

- an affine function $f$ satisfies the superposition

$$
f(\alpha \boldsymbol{x}+\beta \boldsymbol{y})=\alpha f(\boldsymbol{x})+\beta f(\boldsymbol{y})
$$

for any affine combination $\alpha+\beta=1$

- $f: \mathbb{R}^{n} \rightarrow \mathbb{R}$ is affine if we can write it as

$$
f(\boldsymbol{x})=\boldsymbol{a}^{T} \boldsymbol{x}+b
$$

for some $n$-vector $\boldsymbol{a}$ and scalar $b$ (linear function plus a constant)

## Quadratic functions

a function $f: \mathbb{R}^{n} \rightarrow \mathbb{R}$ is quadratic if it can be expressed as

$$
f(\boldsymbol{x})=\boldsymbol{x}^{T} Q \boldsymbol{x}+\boldsymbol{x}^{T} \boldsymbol{r}+c
$$

where

- $Q$ is an $n \times n$ matrix
- $r$ is an $n$-vector
- $c$ is a scalar


## Quadratic Form:

- a quadratic form is a special case: $\boldsymbol{x}^{T} Q \boldsymbol{x}$ where $Q$ is symmetric
- we can always assume $Q$ is symmetric because:

$$
\boldsymbol{x}^{T} Q \boldsymbol{x}=(1 / 2) \boldsymbol{x}^{T}\left(Q+Q^{T}\right) \boldsymbol{x}
$$

hence, $\boldsymbol{x}^{T} Q \boldsymbol{x}=\boldsymbol{x}^{T} P \boldsymbol{x}$ with $P=\frac{1}{2}\left(Q+Q^{T}\right)$ being symmetric

## References

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